import java.util.Random;

```
public class Puzzle5 {
```

private static Random rnd = new Random();

```
public static void main(String[] args) {
   StringBuffer word = null;
   switch(rnd.nextInt(2)) {
      case 1: word = new StringBuffer('R');
      case 2: word = new StringBuffer('T');
      default: word = new StringBuffer('J');
}
```

word.append("est"); System.out.println(String.format("It's time for a %s.", word));

#### public class Puzzle6 {

public static void main(String[] args) {
 System.out.println(getResult());

private static boolean getResult() {
 try {
 return true;
 } finally {
 return false;
 }
}

# Ice Breaker Exercise: EIL5 "Programming"

- Explain It Like I'm Five (EIL5): How do computer programs work?
- Can your explanation intuitively address:
  - Complexity of software
  - Programming bugs
  - Security issues



# Control Flow Graph (CFG)

- A control flow graph (CFG) is a graph representation that captures the paths that might be traversed through a program during its execution, (i.e. the orderings that the program's statements may be executed in at runtime).
- Reading: Frances E. Allen. 1970. Control flow analysis. In Proceedings of a symposium on Compiler optimization. ACM, New York, NY, USA, 1-19.

# Control Flow Graph (CFG)







# **Counting Program Paths**

• How many paths are there for *n* nested branches?

```
if(condition_1){
  if(condition 2){
    if(condition_3){
       ...
       if(condition_n){
         // conditions 1 through n
         // must all be true to reach here
```



# **Counting Program Paths**

• How many paths are there for *n* non-nested branches?

```
if(condition_1){
  // code block 1
if(condition_2){
  // code block 2
if(condition_3){
  // code block 3
• • •
if(condition_n){
  // code block n
```

Ļ



# Considering Loops

- Programs may have loops
  - How many paths does this program have?
  - Can we say if this program halts?

```
while(condition_1){
    if(condition_2){
        break;
    }
}
```



# The Halting Problem

Suppose, we could construct: H(M, x) := if M halts on x then return true else return false

Then we could construct: G(M, x) := if G(M, x) is false then return true else loop forever

But if we then pass G to itself, that is G(G,G), we get a contradiction between what G does and what H says that G does. If H says that G halts, then G does not halt. If H says that G does not halt, then it does halt.

H cannot exist.





### Group Formation

• Count off 1 to 4

# Eclipse Plugin Development + Atlas

- Setting Up Eclipse with Atlas
  - Install: Atlas (all plugins), Plugin Development Tools
- Debug As  $\rightarrow$  Eclipse Application
- Xinu Project (included with Atlas)
- Atlas Shell (add plugin project to dependencies)

Push A (root) on stack.

Stack: [A] History: Paths:



Pop A onto history. A is not a leaf so push C,B.

Stack: [C, B] History: [A] Paths:



Pop B onto history. B is a leaf so save path and trim history.

Stack: [C] History: [A, <del>B</del>] Paths: [A, B]



Pop C onto history. C is not a leaf so push E,D.

Stack: [E, D] History: [A, C] Paths: [A, B]



Pop D onto history. D is a leaf so save path and trim history.

Stack: [E] History: [A, C, <del>D</del>] Paths: [A, B], [A, C, D]



Pop E onto history. E is a leaf so save path and trim history.

Stack: [] History: [A, C, <del>E</del>] Paths: [A, B], [A, C, D], [A, C, E]



Stack is empty. Paths are enumerated.

Stack: [] History: [A, C] Paths: [A, B], [A, C, D], [A, C, E]



# Problem 2 (remaining class time)

- Discuss path counting strategy
  - Why do we need a DAG? What's the implications of using a DAG?
  - What is the complexity of the algorithm with respect to the DAG?
- Explore support code
  - CFGPrinter.java Example
- Group Coding for problem 2